## Random Graphs — Assignment 1

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Question 1. Fix a graph H, and suppose that  $p = \omega(n^{-1/m(H)})$ . In class, we showed that with high probability, G(n, p) contains a copy of H. Show that in fact, with high probability G(n, p) contains  $\omega(1)$  copies of H.

Answer. Let X be the number of copies of H in G(n, p). In class we showed that  $\mathbb{E}[X] = \omega(1)$  and  $\mathbb{V}[X] = o(\mathbb{E}[X]^2)$ . Hence Chebyshev's inequality shows that

$$\Pr[X \le \mathbb{E}[X]/2] \le \Pr[|X - \mathbb{E}[X]| \ge \mathbb{E}[X]/2] \le \frac{\mathbb{V}[X]}{\mathbb{E}[X]^2/4} = o(1).$$

It follows that with high probability,  $X \ge \mathbb{E}[X]/2 = \omega(1)$ .

**Question 2.** Suppose that p = o(1/n). Estimate the probability that G(n, p) contains a triangle by finding a function q(n, p) such that

 $\Pr[G(n, p) \text{ contains a triangle}] \sim q(n, p).$ 

(Note that as shown in class, q(n, p) = o(1).)

Answer. Let X be the number of triangles in G(n, p). The Bonferroni inequalities give the bounds

$$\mathbb{E}[X] - \mathbb{E}\left[\binom{X}{2}\right] \le \Pr[X > 0] \le \mathbb{E}[X].$$

In class we have shown that

$$\mathbb{E}[X] \sim (np)^3/6,$$
$$\mathbb{E}\left[\binom{X}{2}\right] = O(n^6p^6 + n^5p^6 + n^4p^5) = \mathbb{E}[X] \cdot O(n^3p^3 + n^2p^3 + np^2) = o(\mathbb{E}[X]).$$

It follows that

$$\Pr[X > 0] \sim \mathbb{E}[X] \sim (np)^3/6.$$

Question 3. In this question we will analyze the probability that G(n, c/n) contains the graph  $H := \triangleright$  (a triangle with an attached edge).

(a) Let  $E_k$  be the event that G(n, c/n) contains exactly k vertex-disjoint triangles, each of them isolated (not connected by an edge to the outside world), and no other triangles. Calculate  $\lim_{n\to\infty} \Pr[E_k]$ , and deduce a lower bound on the probability that G(n, c/n) is H-free (contains no copy of H).

*Proof.* For every specific choice of k vertex-disjoint triangles, the probability of the event is

$$(c/n)^{3k}(1-c/n)^{3k(n-3k)+\binom{k}{2}-3k}(e^{-c^3/6}\pm o(1))\sim e^{-c^3/6}(c/n)^{3k}e^{-3ck}.$$

The number of possible choices is  $\frac{1}{k!} \binom{n}{3} \cdots \binom{n-3(k-1)}{3} \sim \frac{1}{k!} (n^3/6)^k$ , and so

$$\Pr[E_k] \sim e^{-c^3/6} \cdot e^{-3ck} \frac{(c^3/6)^k}{k!}.$$

Let us denote the right-hand side by  $q_k$ . We can calculate

$$q := \sum_{k=0}^{\infty} q_k = e^{-c^3/6} \sum_{k=0}^{\infty} \frac{((ce^{-c})^3/6)^k}{k!} = e^{-c^3/6} \cdot e^{(ce^{-c})^3/6} = e^{-(1-e^{-3c})(c^3/6)}.$$

For every  $\epsilon > 0$ , we can find  $\ell$  such that  $\sum_{k=0}^{\ell} q_{\ell} \ge q - \epsilon$ . Hence for every  $\epsilon > 0$ , the probability that G(n, c/p) is *H*-free is at least  $q - \epsilon - o(1)$ . It follows that the probability that G(n, c/p) is *H*-free is at least q - o(1).  $\Box$ 

(b) Show that for every  $\epsilon > 0$  there is k such that for large enough n, the probability that G(n, c/n) contains more than k triangles is less than  $\epsilon$ .

Answer. We know that for every k, the probability that G(n, c/n) contains more than k triangles tends to

$$1 - e^{-c^3/6} \sum_{\ell=0}^{k} \frac{(c^3/6)^{\ell}}{\ell!}.$$

We can find k so that this is at most  $\epsilon/2$ . Hence for large enough n, the probability that G(n, c/n) contains more than k triangles is less than  $\epsilon$ .

(c) Prove a matching upper bound on the probability that G(n, c/n) is H-free.

Answer. If  $G \sim G(n, c/n)$  is *H*-free then either one of the events  $E_0, \ldots, E_k$  happens, or *G* contains more than *k* triangles. Note that *G* cannot contain two overlapping triangles, since any such configuration contains a copy of *H*. The calculation above shows that

$$\lim_{n \to \infty} \sum_{\ell=0}^{k} \Pr[E_{\ell}] < q,$$

and in particular, for large enough n the sum is at most q. The previous item shows that for every  $\epsilon > 0$ , we can find k such that the probability that G contains more than k triangles is at most  $\epsilon$ . In total, we get that for every  $\epsilon > 0$ , for large enough n the probability that G(n, c/n) is H-free is at most  $q + \epsilon$ . Since this holds for every  $\epsilon > 0$ , it follows that the probability that G(n, c/n) is H-free is at most q + o(1).  $\Box$